# Locally private learning without interaction requires separation

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# Local Differential Privacy (LDP)



#### [KLNRS '08]

 $\epsilon$ -LDP if for every user *i*, message *j* is sent using a local  $\epsilon_{i,j}$ -DP randomizer  $A_{i,j}$  and

$$\sum_{j} \epsilon_{i,j} \le \epsilon$$

### Non-interactive LDP



### PAC learning

**PAC model [Valiant '84]**: Let *C* be a set of binary classifiers over *X A* is a PAC learning algorithm for *C* if  $\forall f \in C$  and distribution *D* over *X*, given i.i.d. examples  $(x_i, f(x_i))$  for  $x_i \sim D$ , *A* outputs *h* such that w.h.p.  $\Pr_{x \sim D}[h(x) \neq f(x)] \leq \alpha$ 

Distribution-specific learning: D is fixed and known to A

# Statistical query model [Kearns '93]



P distribution over Z

 $Z = X \times \{\pm 1\}$ P is the distribution of (x, f(x)) for  $x \sim D$ 

 $\phi_1: Z \to [0,1] |v_1 - \mathbf{E}_{z \sim P}[\phi_1(z)]| \le \tau$  $\tau$  is tolerance of the query;  $\tau = 1/\sqrt{n}$ 

[KLNRS '08] Simulation with success prob.  $1 - \beta$  ( $\epsilon \le 1$ )

- $\epsilon$ -LDP with *m* messages  $\Rightarrow O(m)$  queries with  $\tau = \Omega\left(\frac{\beta}{m}\right)$
- q queries with tolerance  $\tau \Rightarrow \epsilon$ -LDP with  $n = O\left(\frac{q \log(q/\beta)}{(\tau \epsilon)^2}\right)$  samples/messages Non-interactive if and only if queries are non-adaptive

# Known results

C is SQ-learnable efficiently (non-adaptively) if and only if learnable efficiently with  $\epsilon$ -LDP (non-interactively)

Examples:

- Yes: halfspaces/linear classifiers [Dunagan,Vempala '04]
- No: parity functions [Kearns '93]
- Yes, non-adaptively: Boolean conjunctions

[KLNRS 08] There exists C that is

- 1. SQ/LDP-learnable efficiently over the uniform distribution on  $\{0,1\}^d$  but
- 2. requires exponential num. of samples to learn non-interactively by an LDP algorithm



[KLNRS 08]:

Does separation hold for distribution-independent learning?

Masked parity

# Margin Complexity



Margin complexity of *C* over *X* - MC(*C*): smallest *M* such that exists an embedding  $\Psi: X \to \mathbf{B}_d(1)$  under which every  $f \in C$  is linearly separable with margin  $\gamma \ge \frac{1}{M}$ 

Positive examples {  $\Psi(x) | f(x) = +1$ } Negative examples {  $\Psi(x) | f(x) = -1$ }

#### Lower bound

Thm: Let C be a negation-closed set of classifiers. Any non-interactive 1-LPD algorithm that learns C with error  $\alpha < 1/2$  and success probability  $\Omega(1)$  needs  $n = \Omega\left(\mathbf{MC}(C)^{2/3}\right)$ 

Corollaries:

- Decision lists over  $\{0,1\}^d$ :  $n = 2^{\Omega(d^{1/3})}$ [Buhrman,Vereshchagin,de Wolf '07] (Interactively) learnable with  $n = poly\left(\frac{d}{\alpha\epsilon}\right)$  [Kearns '93]
- Linear classifiers over  $\{0,1\}^d$ :  $n = 2^{\Omega(d)}$ [Goldmann,Hastad,Razborov '92; Sherstov '07] (Interactively) learnable with  $n = poly\left(\frac{d}{\alpha\epsilon}\right)$  [Dunagan,Vempala '04]



Thm: For any C and distribution D there exists a non-adaptive  $\epsilon$ -LPD algorithm that learns C over D with error  $\alpha$  and success probability  $1 - \beta$  using

$$n = \operatorname{poly}\left(\operatorname{MC}(C) \cdot \frac{\log(1/\beta)}{\alpha \epsilon}\right)$$

Instead of fixed D

- access to public unlabeled samples from D
- (interactive) LDP access to unlabeled samples from D

Lower bound holds against the hybrid model

Thm: Let C be a negation-closed set of classifiers. If exists a non-adaptive SQ algorithm that uses q queries of tolerance 1/q to learn C with error  $\alpha < 1/2$  and success probability  $\Omega(1)$  then  $MC(C) = O(q^{3/2})$ 

**Correlation dimension of** *C* **- CSQdim(***C***) [F. '08]** : smallest *t* for which exist *t* functions  $h_1, ..., h_t: X \to [-1,1]$  such that for every  $f \in C$  and distribution *D* exists *i* such that  $\left| \sum_{x \sim D} [f(x)h_i(x)] \right| \ge \frac{1}{t}$ 

Thm: [F. '08; Kallweit, Simon '11]:

 $MC(C) \le CSQdim(C)^{3/2}$ 

#### Proof

If exists a non-adaptive SQ algorithm A that uses q queries of tolerance 1/q to learn C with error  $\alpha < 1/2$  then  $CSQdim(C) \le q$ 

Let  $\phi_1, ..., \phi_q: X \times \{\pm 1\} \rightarrow [0,1]$  be the (non-adaptive) queries of A Decompose

$$\phi(x,y) = \frac{\phi(x,1) + \phi(x,-1)}{2} + \frac{\phi(x,1) - \phi(x,-1)}{2} \cdot y$$

 $\mathop{\mathbf{E}}_{x \sim D}[\phi_i(x, f(x))] = \mathop{\mathbf{E}}_{x \sim D}[g_i(x)] + \mathop{\mathbf{E}}_{x \sim D}[f(x)h_i(x)]$ 

If 
$$\left| \underset{x \sim D}{\mathbf{E}} [f(x)h_i(x)] \right| \le \frac{1}{q}$$
 then  $\underset{x \sim D}{\mathbf{E}} [\phi_i(x, f(x))] \approx \underset{x \sim D}{\mathbf{E}} [\phi_i(x, -f(x))]$ 

If this holds for all  $i \in [q]$ , then the algorithm cannot distinguish between f and -fCannot achieve error < 1/2 Thm: For any C and distribution D there exists a non-adaptive  $\epsilon$ -LPD algorithm that learns C over D with error  $\alpha < 1/2$  and success probability  $1 - \beta$  using  $\log(1/\beta)$ 

$$n = \operatorname{poly}\left(\operatorname{MC}(C) \cdot \frac{\operatorname{log}(1/p)}{\alpha \epsilon}\right)$$

Margin complexity of *C* over *X* - MC(*C*): smallest *M* such that exists an embedding  $\Psi: X \to \mathbf{B}_d(1)$  under which every  $f \in C$  is linearly separable with margin  $\gamma \ge \frac{1}{M}$ 

Thm [Arriaga,Vempala '99; Ben-David,Eiron,Simon '02]: For every every  $f \in C$ , random projection into  $\mathbf{B}_d(1)$  for  $d = O(\mathbf{MC}(C)^2 \log(1/\beta))$ ensures that with prob.  $1 - \beta$ ,  $1 - \beta$  fraction of points are linearly separable with margin  $\gamma \ge \frac{1}{2 \operatorname{MC}(C)}$ 

# Algorithm

Perceptron: if sign( $\langle w_t, x \rangle$ )  $\neq y$  then update  $w_{t+1} \leftarrow w_t + yx$ Expected update:  $\underset{(x,y)\sim P}{\mathbf{E}}[yx \mid \text{sign}(\langle w_t, x \rangle) \neq y]$  $\begin{bmatrix} \mathbf{E}\\(x,y)\sim P \begin{bmatrix} yx \cdot \mathbb{1}(\text{sign}(\langle w_t, x \rangle) \neq y) \end{bmatrix} / \underbrace{\Pr_{(x,y)\sim P}[\text{sign}(\langle w_t, x \rangle) \neq y]}_{(x,y)\sim P} \text{ scalar} \geq \alpha \\ \end{bmatrix} \text{ scalar} \geq \alpha \\ \begin{bmatrix} \mathbf{E}\\(x,y)\sim P \begin{bmatrix} x \cdot \frac{y - \text{sign}(\langle w_t, x \rangle)}{2} \end{bmatrix} = \underbrace{\frac{\mathbf{E}\\(x,y)\sim P}[xy] + \underbrace{\mathbf{E}}\\(x,y)\sim P}[x \text{ sign}(\langle w_t, x \rangle)]}_{\text{non-adaptive}} \text{ independent of the label}$ 

Estimate the mean vector with  $\ell_2$  error

- LDP [Duchi,Jordan,Wainright '13]
- SQs [F.,Guzman,Vempala '15]

#### Conclusions

- New approach to lower bounds for non-interactive LDP

   Reduction to margin-complexity lower bounds
- Lower bounds for classical learning problems
- Same results for communication constrained protocols
   Also equivalent to SQ
- Interaction is necessary for learning
- Open:
  - Distribution-independent learning in poly(MC(C))
  - Lower bounds against 2 + round protocols
  - Stochastic convex optimization

https://arxiv.org/abs/1809.09165

